Chair of Symbolic Computation Prof. Dr. Martin Kreuzer Julian Danner



## Computeralgebra – Sheet 1

Summer term 2025

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Exercise 1 (Extended Euclidean Algorithm).

Let  $a, b \in \mathbb{Z}$ . Consider the following sequence of instructions.

- (1) If a = 0 and b = 0, **return** the triple [0, 0, 0]. If a = 0 and  $b \neq 0$ , **return** the triple  $[0, \frac{|b|}{b}, |b|]$ . If  $a \neq 0$  and b = 0, **return** the triple  $[\frac{|a|}{a}, 0, |a|]$ .
- (2) Set the triples  $[c_0, d_0, e_0] \leftarrow [\frac{|a|}{a}, 0, |a|]$  and  $[c_1, d_1, e_1] \leftarrow [0, \frac{|b|}{b}, |b|].$
- (3) If  $e_0 < e_1$ , swap  $[c_0, d_0, e_0] \longleftrightarrow [c_1, d_1, e_1]$ .
- (4) Repeat steps (4.1)–(4.3) until  $e_1 = 0$ .
  - (4.1) Write  $e_0$  in the form  $e_0 = qe_1 + r$ , where  $q, r \in \mathbb{N}$  and  $0 \le r < e_1$ .
  - (4.2) Compute  $[c_2, d_2, e_2] \leftarrow [c_0 qc_1, d_0 qd_1, r].$
  - (4.3) Assign  $[c_0, d_0, e_0] \leftarrow [c_1, d_1, e_1]$  and  $[c_1, d_1, e_1] \leftarrow [c_2, d_2, e_2]$ .
- (6) **return** the triple  $[c_0, d_0, e_0]$ .

This algorithm is known as the *Extended Euclidean Algorithm* (EEA) and computes  $[c, d, e] \in \mathbb{Z}^3$  such that  $e = \gcd(a, b)$  and ac + bd = e.

- (a) Show that this is indeed an algorithm, i.e., that it stops after finitely many steps.
- (b) Show that the output is correct, i.e., a triple with the claimed properties.
  - *Hint:* Show that  $ac_0 + bd_0 = e_0$  is an invariant of the algorithm, then it suffices to show that e = gcd(a, b).

**Exercise 2.** Let p be a prime and let  $\mathbb{F}_p = \mathbb{Z}/p\mathbb{Z}$  be the finite field with p elements. Explain how the *Extended Euclidean Algorithm* can be used to compute the inverse of a unit in  $\mathbb{F}_p$ .

**Exercise 3.** Let K be a field, P = K[x], let  $f_1, f_2 \in P$  be non-zero polynomials, and let

$$I = Pf_1 + Pf_2 = \{ g_1f_1 + g_2f_2 \mid g_1, g_2 \in P \}.$$

- (a) Show that the set I is an ideal of P.
- (b) Show that the ideal I is generated by gcd(f<sub>1</sub>, f<sub>2</sub>), i.e., I = ⟨gcd(f<sub>1</sub>, f<sub>2</sub>)⟩.
  Recall that the greatest common divisor of f<sub>1</sub> and f<sub>2</sub>, denoted by gcd(f<sub>1</sub>, f<sub>2</sub>), is the unique monic polynomial h ∈ P with h | f<sub>1</sub>, h | f<sub>2</sub>, and if g | f<sub>1</sub> and g | f<sub>2</sub> then g | h.

**Exercise 4.** Let p be a prime number, let  $K = \mathbb{F}_p$ , let  $f \in P = K[x]$  be a non-zero polynomial, and let f' be the derivative of f. Show that f' = 0 if and only if f is of the form  $f = g^p$  for some  $g \in P$ .

**Exercise 5.** Let p be a prime number, let  $K = \mathbb{F}_p$  or  $\mathbb{Q}$ , let  $f \in K[x]$  be a non-zero polynomial, and let f' be the derivative of f. Show that if f is irreducible then we have gcd(f, f') = 1.